

Language Design and Overview of COOL

CS143

Lecture 2

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Grade Weights

- Project 50%
 - 1–2 10% each
 - 3–4 15% each
- Midterm 15%
- Final 25%
- Written Assignments 10%
 - 2.5% each

Lecture Outline

- Today's topic: language design
 - Why are there new languages?
 - Good-language criteria
- History of ideas:
 - Abstraction
 - Types
 - Reuse
- Cool
 - The Course Project

Programming Language Economics 101

- Languages are adopted to fill a void
 - Enable a previously difficult/impossible application
 - Orthogonal to language design quality (almost)
- Programmer training is the dominant cost
 - And rewriting code
 - Languages with many users are replaced rarely
 - Popular languages become ossified
 - But easy to start in a new niche . . .

Why So Many Languages?

- Application domains have distinctive and conflicting needs
- Examples: (language–need pairs)

Topic: Language Design

- No universally accepted metrics for design
- Claim: “A good language is one people use”

Language Evaluation Criteria

Features	Criteria		
	Readability	Writeability	Reliability
Data types			
Abstraction			
Type checking			
Exception handling			

History of Ideas: Abstraction

- Abstraction = detached from concrete details
 - “Abstraction is selective ignorance” - Andrew Koenig
- Abstraction is necessary to build any complex system
 - The key is **information hiding**—expose only the essential
- Modes of abstraction
 - **Via languages/compilers**: High-level code, few machine dependencies
 - **Via functions and subroutines**: Abstract interface to behavior
 - **Via modules**: Export interfaces; hide implementation
 - **Via classes/abstract data types**: Bundle data with its operations

History of Ideas: Types

- Originally, few types
 - FORTRAN: scalars, arrays
 - LISP: no static type distinctions
- Realization: Types help
 - Lets you to express abstraction
 - Lets the compiler report many frequent errors
 - Sometimes to the point that programs are guaranteed “safe”
 - Helps the compiler optimize your code
- More recently
 - Lots of interest in types
 - Experiments with various forms of parameterization
 - Best developed in functional programming

History of Ideas: Reuse

- Reuse = exploit common patterns in software systems
 - Goal: mass-produced software components
 - Reuse is difficult
- Two popular approaches
 - Type parameterization (`List(int)`, `List(double)`)
 - Classes and inheritance: C++ derived classes
 - C++ and Java have both
- Inheritance allows
 - Specialization of existing abstraction
 - Extension, modification, and hidden behavior

Trends

- Language design
 - Many new special-purpose languages
 - Popular languages stick around (perhaps forever)
 - Fortran and Cobol
- Compilers
 - Ever more needed and ever more complex
 - Driven by increasing gap between
 - new languages
 - new architectures
 - Venerable and healthy area

Why Study Languages and Compilers ?

5. Increase capacity of expression
4. Improve understanding of program behavior
3. Increase ability to learn new languages
2. Learn to build a large and reliable system
1. See many basic CS concepts at work

Cool Overview

- Classroom Object Oriented Language
- Designed to
 - Be implementable in a short time
 - Give a taste of implementation of modern
 - Abstraction
 - Static typing
 - Reuse (inheritance)
 - Memory management
 - And more ...
- But many things are left out

A Simple Example

```
class Point {  
    x : Int ← 0;  
    y : Int ← 0;  
};
```

- Cool programs are sets of class definitions
 - A special class **Main** with a special method **main**
 - All Cool code lives inside classes
- A class is a collection of attributes and methods
- Instances of a class are objects

Cool Objects

```
class Point {  
    x : Int ← 0;  
    y : Int; (* use default value *)  
};
```

- The expression “**new Point**” creates a new object of class **Point**
- An object can be thought of as a record with a slot for each attribute

x	y
0	0

Methods

- A class can also define methods for manipulating the attributes

```
class Point {
  x : Int ← 0;
  y : Int ← 0;
  movePoint(newx : Int, newy : Int): Point {
    {
      x ← newx;
      y ← newy;
      self;
    } -- close block expression
  }; -- close method
}; -- close class
```

- Methods can refer to the current object using `self`

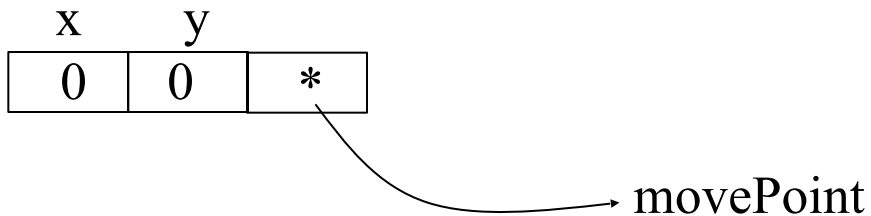
Information Hiding in Cool

- Methods are global
- Attributes are local to a class
 - They can only be accessed by the class's methods

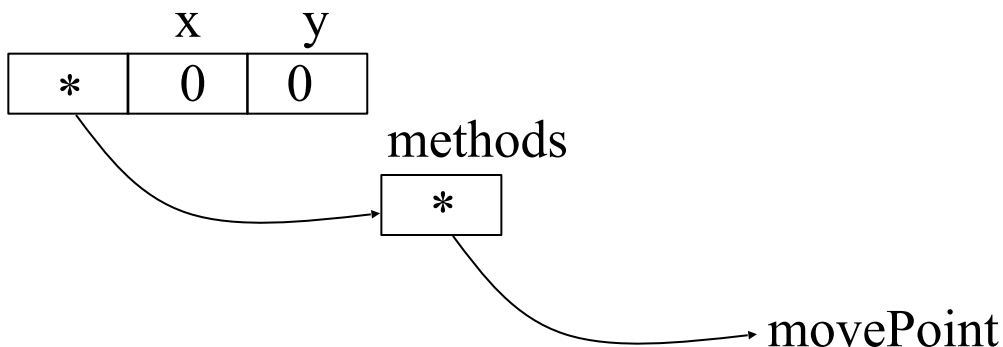
```
class Point {  
    ...  
    x () : Int { x };  
    setx (newx : Int) : Int { x ← newx };  
};
```

Methods

- Each object knows how to access the code of a method
 - As if the object contains a slot pointing to the code



- In reality implementations save space by sharing these pointers among instances of the same class



Inheritance

- We can extend points to colored points using subclassing
=> class hierarchy

```
class ColorPoint inherits Point {  
  color : Int ← 0;  
  movePoint(newx : Int, newy : Int): Point {{  
    color ← 0;  
    x ← newx;  
    y ← newy;  
    self;  
  }};  
};
```

movePoint	x	y	color
*	0	0	0

*	0	0	0
---	---	---	---

Cool Types

- Every class is a type
- Base classes:
 - `Int` for integers
 - `Bool` for boolean values: `true`, `false`
 - `String` for strings
 - `Object` root of the class hierarchy
- All variables must be declared
 - compiler infers types for expressions

Cool Type Checking

```
x : A;  
x ← new B;
```

- Is well typed if **A** is an ancestor of **B** in the class hierarchy
 - Anywhere an **A** is expected a **B** can be used
- Type safety:
 - A well-typed program cannot result in runtime type errors

Method Invocation and Inheritance

- Methods are invoked by dispatch
- Understanding dispatch in the presence of inheritance is a subtle aspect of OO languages

```
p : Point;
p ← new ColorPoint;
p.movePoint(1,2);
```

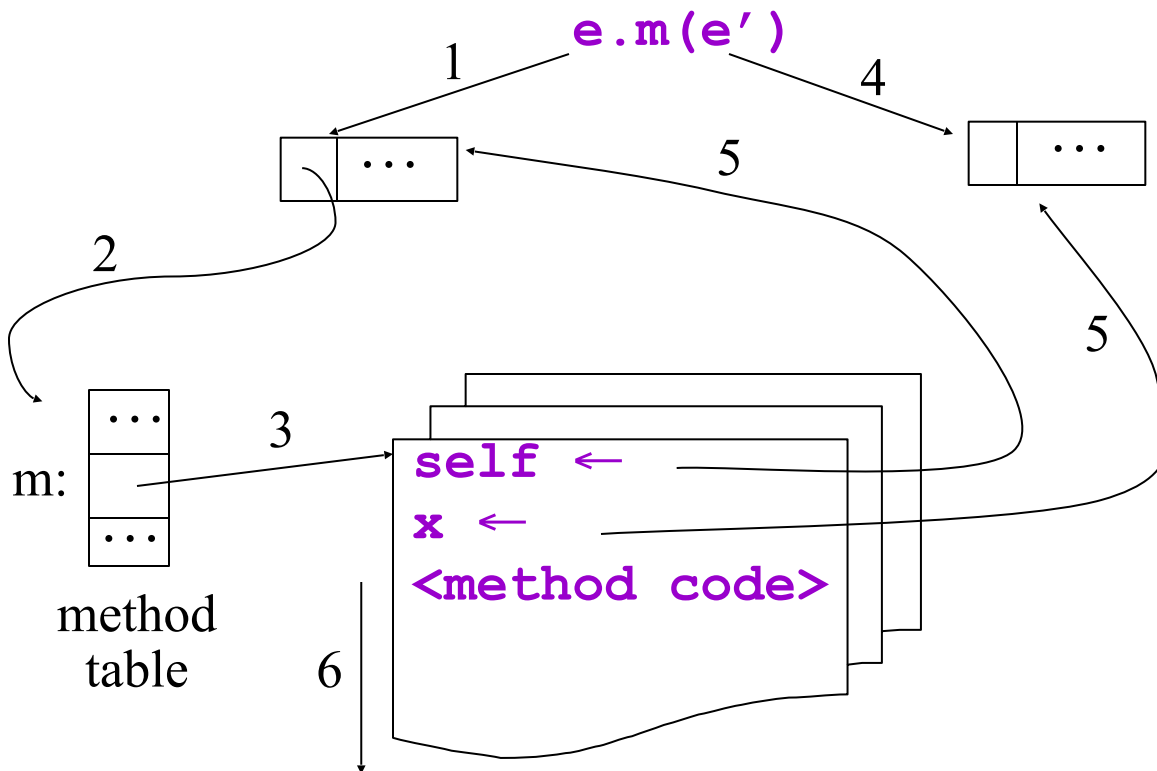
← p has static type **Point**

← p has dynamic type **ColorPoint**

← p.movePoint must invoke the **ColorPoint** version

Method Invocation

- Example: invoke one-argument method $m(x)$



1. Eval. e
2. Find class of e
3. Find code of m
4. Eval. argum.
5. Bind `self` and x
6. Run method

Other Expressions

- Expression language
 - every expression has a type and a value
 - Loops `while E loop E pool`
 - Conditionals `if E then E else E fi`
 - Case statement `case E of x : Type ⇒ E; ... esac`
 - Arithmetic `+, -, ...`
 - Logical operations `<, =, ...`
 - Assignment `x ← E`
 - Primitive I/O `out_string(s), in_string(), ...`
- Missing features:
 - arrays, floating point operations, exceptions, ...

Cool Memory Management

- Memory is allocated every time `new` is invoked
- Memory is deallocated automatically when an object is no longer reachable
- Done by the garbage collector (GC)
 - There is a Cool GC

Course Project

- A complete compiler
 - Cool ==> MIPS assembly language
 - No optimizations
- Split in 4 programming assignments (PAs)
- There is adequate time to complete assignments
 - But start early and please follow directions
- Individual or team
 - max. 2 students